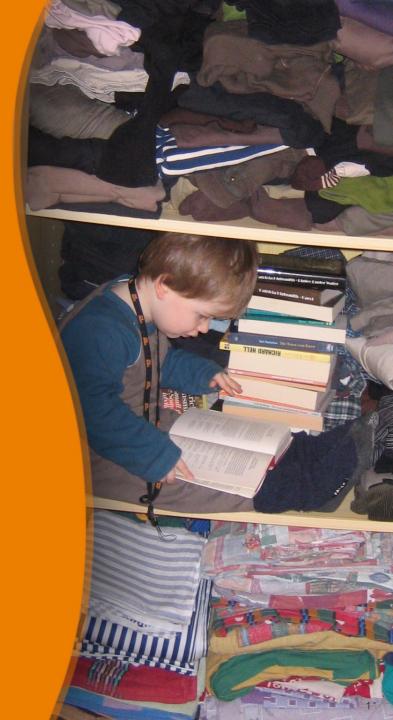


UNO: Anecdotal Evidence

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If art is the tip of the iceberg I'm the part sinking below —Lou Reed, John Cale, Songs for Drella



UNO

- An object-oriented framework to bring together:
 - > different programming languages (C++, Java, Python, ...)
 - > different environments
 - within a process (C++ runtime, JVM, ...)
 - across processes (named pipes, sockets)
 - across machines (sockets)
- Using bridges among the environments to marshal method invocations on objects as remote procedure calls.
 - Where the shortcomings of this approach are well known.



Breaking no new ground —Bongwater



Minimal Type System

- UNO type system reflects capabilities of target languages:
 - void, boolean, char and string, various integral and floating point types (minimal?)
 - > enums (modeled after C)
 - > structs (records), supporting parametric polymorphism
 - > interfaces (references to objects, incl. null)
 - > sequences
 - > exceptions (somewhat second class)
 - > any (type + value)



No Algebraic Data Types, etc.

- Overall, a rather mediocre type system:
 - > Encoding sums with products,
 struct Optional<T>{boolean isPresent; T Value;};
 instead of
 Optional<T> = NotPresent | Value T
 - No parametric polymorphism for interfaces: XInterface getElement(); or any getElement();
- Then again, how to map the good concepts to the relevant target languages?



Java Binding Oddities

- The base UNO XInterface is mapped to java.lang.Object. But UNO objects still have to implement XInterface.
- UNO any is mapped to java.lang.Object, too:

long

→ java.lang.Integer or com.sun.star.uno.Any

unsigned long

→ com.sun.star.uno.Any

XInterface reference

→ java.lang.Object or com.sun.star.uno.Any

XFoo reference

→ com.sun.star.uno.Any



Identifiers

- Naively modeled after C.
- Problems with name clashes in language bindings (keywords, methods named the same as containing interfaces, etc.).
- Restrictions retrofitted to allow escape mechanisms in language bindings (goto → method_goto).
- What about languages where case has meaning?



Object Life-Cycle

- Implemented with reference counting.
 - > C++ Reference<T> calls acquire/release.
 - Java finalize calls release across bridge.
- C++ programmers confused when their destructors are called late (only after JVM garbage collection).
- Cyclic references cause problems, of course.
 - > XComponent.dispose
 - > XCloseBroadcaster, XCloseable, and CloseVetoException



URP Oddities

- Compact wire representation.
- Marshalling types:

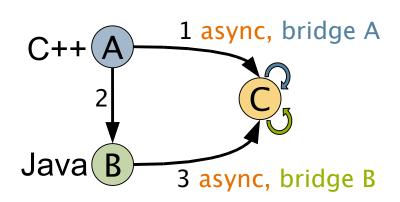
```
long: 0x06
com.sun.star.uno.XInterface: 0x16 "com.sun.st..."
sequence<long>: 0x14 "[]long"
instead of 0x14 0x06
```

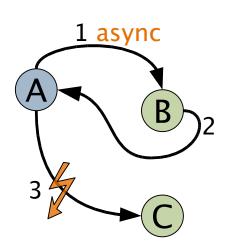
- Negotiating protocol properties during startup:
 - If both sides want to set the same properties, both send the same marshalled blob plus a random number. Unnecessary life-lock if both sides use identical random numbers.



Asynchronous One-Way Calls

- UNO methods marked as [oneway], implemented as (semi-)asynchronous URP calls.
 - > All async calls executed in order.
 - > All async calls executed before next synchronous call.
 - > For each UNO thread, there is at most one concurrent thread executing the async calls.
- Does not really work:







Evolution

- User-defined UNO types want to evolve over time.
- This causes problems in two dimensions:
 - Interfaces have both clients and implementations.
 - Implementation languages offer limited support.
- UNO only offers mechanisms, no policies.
 - > Evolve XFoo to XFoo2 to XFoo3.
 - > Published vs. unpublished types.



Everyone wants the honey but not the sting

—Mark Edwards



Runtime Type Information

- UNO types described in textual .idl files.
- Translated and combined into binary .rdb files.
- Read at runtime:
 - Packing values (type description + void*) in anys.
 - Marshalling data across bridges.
 - Creating vtables of C++ proxy objects.
- For Java, .class files are used instead.
- So there are two (three, even) descriptions for each type.



.idl Files

- Are C-style preprocessed, complicating things without need.
 - > Preprocessor used to handle includes.
 - > But code makers have to be careful to not generate code for merely included types.
 - Could have used the Java way instead, as (almost) each type is in its own .idl file, anyway.



.rdb Files

- Back when StarOffice still had mail/news functionality, it used storages to store mails in. Those often broke and caused data loss.
- A new implementation solved the stability issues.
 - But it is unsound, reducing arbitrary-size string keys to fixed length hash values.)
- This was re-used for the UNO type database.
 - Complex code (no direct access from Java).
 - > Files are unnecessarily large, due to lots of padding.



Active Component Registration

- Each UNO component (collection of UNO service implementations; shared library or .jar) must be registered in a services.rdb.
- This is done by calling code in the component.
- OOo extensions are uninstalled as follows:
 - Call code in extension to register its services into a temporary .rdb.
 - Iterate entries in temporary .rdb and remove them from the services.rdb.
 - What if extension code no longer loads after an OOo upgrade?



Dynamically Generated Code

- cppumaker only generates headers, no code.
- C++ proxy objects with vtables are generated on the fly.
- The relevant code is compiler and platform/CPU dependent.
 - > Probably the biggest obstacle for porters.
 - Subtle bugs, plaguing us for years.
- GCC exception handling needs unique RTTI.
 - Implemented via weak symbols, which are expensive at library load time.



Q?

throw SQLException(...);
 ... UNO bridge in between ...
 catch (SQLException e) ...

kept mysteriously failing under Solaris (catch (...) worked).





Q? Ah!

throw SQLException(...);
 ... UNO bridge in between ...
 catch (SQLException e) ...

kept mysteriously failing under Solaris (catch (...) worked).

- In the Solaris C++ ABI, mangled names use the letter Q as an escape mechanism.
- The Solaris C++ bridge ignored that detail when generating mangled names for exception handling.



Instruction Manual Details

On Mac OS X PowerPC, is it

dcbf 0, p or icbi 0, p sync isync

dcbst 0, p sync icbi 0, p isync

?!?





Instruction Manual Details

On Mac OS X PowerPC, is it

```
dcbr 0, p or dcbst 0, p sync icbi 0, p icbi 0, p icbi 0, p icbi 1, p icbi 1, p icbi 2, p icbi 2, p icbi 2, p icbi 2, p icbi 3, p icbi 2, p icbi 2, p icbi 3, p icbi 4, p icbi 6, p icbi 6, p icbi 7, p icbi 1, p icbi 1,
```

 The difference is random crashes when the C++ bridge dynamically generates code and does not correctly synchronize data and instruction memory.



Security-Enhanced Linux

Dynamically generating code under SELinux:
 Instead of

```
p = mmap(0, n, PROT_READ|PROT_WRITE,
   MAP PRIVATE MAP ANON, -1, 0;
 mprotect(p, n,
   PROT READ | PROT WRITE | PROT EXEC);
one now needs
 fd = mkstemp(name);
 unlink(name);
 ftruncate(fd, n);
 p1 = mmap(0, n, PROT READ|PROT WRITE,
   MAP SHARED, fd, 0);
 p2 = mmap(0, n, PROT_READ|PROT_EXEC,
   MAP SHARED, fd, 0);
```



French music sucks but it is nice

—Sonny Vincent

